

Exercise S5.1 – 3-colorable graphs

In the lecture you have seen an algorithm that finds an $O(\sqrt{|V|})$ -coloring for every 3-colorable graph (V, E) . Fix $\alpha = \sqrt{|V|}$. The algorithm goes as follows.

1. While there is a vertex v with degree at least α , color v with a new color, color the neighbors of v with two additional new colors, and delete v and all its neighbors from the graph.
2. Color all remaining vertices with at most $\alpha + 1$ colors.

We want to analyze and generalize this algorithm.

- (a) Where in the algorithm do we use that the graph is 3-colorable?
- (b) How many colors do we need at most (depending on α)? Show that for $\alpha = \sqrt{|V|}$ we only need $O(\sqrt{|V|})$ colors.
- (c) Show that your bound in (a) is tight. More precisely, create a 3-colorable graph on which the algorithm uses $\Omega(\sqrt{|V|})$ colors.

- (d) Can you describe an algorithm that finds a $O\left(|V|^{\frac{2}{3}}\right)$ -coloring for every 4 colorable graph?
- (e) Can you describe an algorithm that finds a $O\left(|V|^{\frac{q-2}{q-1}}\right)$ -coloring for every q colorable graph, where $q \geq 2$ is a constant?